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SECURITY

[RBS-2019-004](#)

PS_NTS ActiveX Control ps_ntscrypto.dll
Relative Distinguished Names Handling Buffer Overflows

Vulnerable Program Details

Details for tested products and versions:

Vendor: SG Co Ltd.
Product: PS_NTS ActiveX Control (PS_NTS.dll)
Version: 4.1.1.4
Component: ps_ntscrypto.dll
Component Version: 4.1.1.3

NOTE: Other versions than the one listed above are likely affected.

Credits

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Impact

The PS_NTS ActiveX control bundles the ps_ntscrypto.dll library, which contains various buffer overflows that may allow an attacker to compromise a user's system.

Vulnerability Details

The ps_ntscrypto.dll library bundled with the PS_NTS ActiveX control exports various crypto functionality, where some of the functions of particular interest to these vulnerabilities are e.g. ps_req_auth_string() and ps_get_session_key(). The exported functions can be reached via methods provided by the ActiveX control such as PsWebClientReqAuth() and PsWebGetSKey().

The PsWebClientReqAuth() method for authenticating a client is defined as follows:

```
[id(0x0000000e), helpstring("method PsWebClientReqAuth")]  
BSTR PsWebClientReqAuth(  
    [in] BSTR inCname,  
    [in] BSTR inPasswd,  
    [in] BSTR inRdname);
```

The PsWebGetSKey() method for getting a session key is defined as follows:

```
[id(0x00000028), helpstring("method PsWebGetSKey")]  
BSTR PsWebGetSKey([in] BSTR RDNAME);
```

In both cases, the argument of interest is 'inRdname' or 'RDNAME', which is intended to supply a list of relative distinguished names.

When either of these two methods is called, the function in PS_NTS.dll that is responsible for handling the call eventually transfers control to the corresponding exported function in ps_ntscrypto.dll along with the supplied arguments. Both functions end up calling a function to parse the supplied relative distinguished name string. This function allocates a 201 byte buffer on the heap and checks if the supplied string starts with "cn=".

```
.text:022D5FA0 ; int __cdecl ConstructRDName(char *pszinRdname)  
.text:022D5FA0 ConstructRDName proc near ; CODE XREF: sub_22D1CD0+53p  
.text:022D5FA0 ; sub_22D1D40+53p ...  
.text:022D5FA0  
.text:022D5FA0 iCount = dword ptr -978h  
.text:022D5FA0 szC = byte ptr -974h ; char[4]  
.text:022D5FA0 iUnk1 = dword ptr -970h  
.text:022D5FA0 wUnk2 = word ptr -96Ch  
.text:022D5FA0 lpszDest = dword ptr -968h  
.text:022D5FA0 szOU = byte ptr -964h ; char[48]  
.text:022D5FA0 anonymous_0 = word ptr -934h  
.text:022D5FA0 szO = byte ptr -930h ; char[52]  
.text:022D5FA0 szCN = byte ptr -8FCh ; char[100]  
.text:022D5FA0 var_898 = byte ptr -898h ; char[200]  
.text:022D5FA0 var_7D0 = byte ptr -7D0h ; char[2000]  
.text:022D5FA0 pszinRdname = dword ptr 4  
.text:022D5FA0  
.text:022D5FA0 sub esp, 978h  
.text:022D5FA6 push ebx  
.text:022D5FA7 push ebp  
.text:022D5FA8 push esi  
.text:022D5FA9 push edi  
.text:022D5FAA mov ecx, 19h  
.text:022D5FAF xor eax, eax  
.text:022D5FB1 lea edi, [esp+988h+szCN] ; char[100]  
.text:022D5FB8 mov ebx, [esp+988h+pszinRdname]  
.text:022D5FBF rep stosd  
.text:022D5FC1 mov dword ptr [esp+988h+szC], eax ; char[4]  
.text:022D5FC5 mov ecx, 0Ch  
.text:022D5FCA mov [esp+988h+iUnk1], eax  
.text:022D5FCE lea edi, [esp+988h+szO] ; char[52]  
.text:022D5FD2 mov [esp+988h+wUnk2], ax  
.text:022D5FD7 rep stosd  
.text:022D5FD9 stosw  
.text:022D5FDB mov ecx, 0Ch  
.text:022D5FE0 xor eax, eax  
.text:022D5FE2 lea edi, [esp+988h+szOU] ; char[48]  
.text:022D5FE6 rep stosd  
.text:022D5FE8 stosw
```

```

.text:022D5FEA          mov     ecx, 32h
.text:022D5FEF          xor     eax, eax
.text:022D5FF1          lea    edi, [esp+988h+var_898] ; char[200]
.text:022D5FF8          test   ebx, ebx
.text:022D5FFA          rep    stosd
.text:022D5FFC          jnz    short loc_22D6009 ; inRdname supplied?
.text:022D6009 loc_22D6009:          ; CODE XREF: ConstructRDName+5Cj
.text:022D6009          push   0C9h ; size_t
.text:022D600E          call   _malloc
.text:022D6013          mov    ebp, eax
.text:022D6015          mov    ecx, 32h
.text:022D601A          xor    eax, eax
.text:022D601C          mov    edi, ebp
.text:022D601E          rep    stosd
.text:022D6020          push   3 ; size_t
.text:022D6022          push   offset aCn_0 ; "cn="
.text:022D6027          push   ebx ; char *
.text:022D6028          mov    [esp+998h+lpszDest], ebp
.text:022D602C          stosb
.text:022D602D          call   _strncmp
.text:022D6032          add    esp, 10h
.text:022D6035          test   eax, eax
.text:022D6037          jnz    short loc_22D6063 ; rdname doesn't start with "cn="?

```

If so, the string is copied to the previously allocated heap buffer without performing any boundary checks. This may cause a heap-based buffer overflow.

```

.text:022D6039          mov    edi, ebx
.text:022D603B          or     ecx, 0FFFFFFFFh
.text:022D603E          repne scasb
.text:022D6040          not    ecx
.text:022D6042          sub    edi, ecx
.text:022D6044          mov    eax, ebp
.text:022D6046          mov    edx, ecx
.text:022D6048          mov    esi, edi
.text:022D604A          mov    edi, ebp
.text:022D604C          shr    ecx, 2
.text:022D604F          rep movsd ; b0f!
.text:022D6051          mov    ecx, edx
.text:022D6053          and    ecx, 3
.text:022D6056          rep movsb
.text:022D6058          pop    edi
.text:022D6059          pop    esi
.text:022D605A          pop    ebp
.text:022D605B          pop    ebx
.text:022D605C          add    esp, 978h
.text:022D6062          retn

```

If, however, the string did not start with “cn=” the function checks if “/C=” can be found anywhere in the supplied string.

```

.text:022D6063 loc_22D6063:          ; CODE XREF: ConstructRDName+97j
.text:022D6063          push   offset aC_0 ; "/C="
.text:022D6068          push   ebx ; char *

```

```
.text:022D6069      call   _strstr
.text:022D606E      mov    edi, eax
.text:022D6070      add    esp, 8
.text:022D6073      test   edi, edi
.text:022D6075      jz     loc_22D636E      ; rdname doesn't contain "/C="?
```

If “/C=” cannot be found, the string is copied into the heap buffer without performing any boundary checks. This may also lead to a heap-based buffer overflow.

```
.text:022D636E loc_22D636E:                                ; CODE XREF: ConstructRDName+D5j
.text:022D636E      mov    edi, ebx
.text:022D6370      or     ecx, 0FFFFFFFh
.text:022D6373      xor    eax, eax
.text:022D6375      repne scasb
.text:022D6377      not    ecx
.text:022D6379      sub    edi, ecx
.text:022D637B      mov    eax, ecx
.text:022D637D      mov    esi, edi
.text:022D637F      mov    edi, ebp
.text:022D6381      shr    ecx, 2
.text:022D6384      rep movsd      ; b0f!
.text:022D6386      mov    ecx, eax
.text:022D6388      mov    eax, ebp
.text:022D638A      and    ecx, 3
.text:022D638D      rep movsb
.text:022D638F      pop    edi
.text:022D6390      pop    esi
.text:022D6391      pop    ebp
.text:022D6392      pop    ebx
.text:022D6393      add    esp, 978h
.text:022D6399      retn
```

If the string does contain “/C=”, a check is performed to see if another “/” character exists later in the string.

```
.text:022D607B      add    edi, 3
.text:022D607E      push   '/'      ; int
.text:022D6080      push   edi      ; char *
.text:022D6081      call   _strchr
.text:022D6086      add    esp, 8
.text:022D6089      test   eax, eax
.text:022D608B      jnz    short loc_22D60AE ; "/" found after "/C="?
```

If so, the whole argument passed to “/C=” is copied into a 4 byte stack buffer without performing any boundary checks. This may lead to a stack-based buffer overflow.

```
.text:022D608D      or     ecx, 0FFFFFFFh
.text:022D6090      lea   edx, [esp+988h+szC] ; char[4]
.text:022D6094      repne scasb
.text:022D6096      not    ecx
.text:022D6098      sub    edi, ecx
.text:022D609A      mov    eax, ecx
.text:022D609C      mov    esi, edi
```

```
.text:022D609E      mov     edi, edx
.text:022D60A0      shr     ecx, 2
.text:022D60A3      rep movsd      ; b0f!
.text:022D60A5      mov     ecx, eax
.text:022D60A7      and     ecx, 3
.text:022D60AA      rep movsb
.text:022D60AC      jmp     short loc_22D60BF
```

If the string does contain a “/” character after “/C=”, the length of the argument to that parameter is calculated before calling `strncpy()` to copy the argument into the 4 byte stack buffer. As the length of the source string is used as size argument, a stack-based buffer overflow may also occur here.

```
.text:022D60AE loc_22D60AE:      ; CODE XREF: ConstructRDName+EBj
.text:022D60AE      sub     eax, edi
.text:022D60B0      lea    ecx, [esp+988h+szC] ; char[4]
.text:022D60B4      push   eax      ; size_t
.text:022D60B5      push   edi      ; char *
.text:022D60B6      push   ecx      ; char *
.text:022D60B7      call  _strncpy  ; b0f!
```

Similar flawed copy operations are later performed for “/O=”, “/OU=”, and “/CN=” when present in the string. These may all result in stack-based buffer overflows.

Solution

The vendor has deprecated the ActiveX control, and KrCERT/CC plans to set the kill-bit.

References

RBS: RBS-2019-004¹
VulnDB: 202041, 202042, 202043, 202044, 202045, 202046

Timeline

2019-01-29	Vulnerability discovered.
2019-02-01	Vulnerability reported to KrCERT/CC.
2019-04-04	Alerts published to VulnDB customers.
2019-05-21	Publication of this vulnerability report.

¹ <https://www.riskbasedsecurity.com/research/RBS-2019-004.pdf>

About Risk Based Security

Risk Based Security offers clients fully integrated security solutions, combining real-time vulnerability and threat data, as well as the analytical resources to understand the implications of the data, resulting in not just security, but the right security.

Company History

Risk Based Security, Inc. (RBS) was established to support organizations with the technology to turn security data into actionable information and a competitive advantage. We do so by enhancing the research available and providing a first of its kind risk identification and evidence-based security management service.

As a data driven and vendor neutral organization, RBS is able to deliver focused security solutions that are timely, cost effective, and built to address the specific threats and vulnerabilities most relevant to the organizations we serve. We not only maintain vulnerability and data breach databases, we also use this information to inform our entire practice.

Solutions

VulnDB - Vulnerability intelligence, alerting, and third party library tracking based on the largest and most comprehensive vulnerability database in the world. Available as feature-rich SaaS portal or powerful API. Vendor evaluations including our Vulnerability Timeline and Exposure Metrics (VTEM), Cost of Ownership ratings, Code Maturity, and Social Risk Scores.

Cyber Risk Analytics - Extensive data breach database including interactive dashboards and breach analytics. Clients are able to gather and analyze security threat and data breach information on businesses, industries, geographies, and causes of loss. It also allows monitoring of domains for data breaches and leaked credentials as well as implementing a continuous vendor management program with our PreBreach data.

YourCISO - Revolutionary service that provides organizations an affordable security solution including policies, vulnerability scans, awareness material, incident response, and access to high quality information security resources and consulting services.

Vulnerability Assessments (VA) and Pentesting - Regularly scheduled VAs and pentests help an organization identify weaknesses before the bad guys do. Managing the most comprehensive VDB puts us in a unique position to offer comprehensive assessments, combining the latest in scanning technology and our own data. Detailed and actionable reports are provided in a clear and easy to understand language.